## Dilip Menariya, D. B. Ojha / International Journal of Engineering Research and Applications (IJERA) ISSN: 2248-9622 www.ijera.com Vol. 2, Issue 5, September- October 2012, pp.328-332 A vital application of security with biometric templates

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#### Abstract

In this Paper, We will discuss high level of biometric templates security is critical issue in personal base authentication system due to addition of error in transmission channel. Therefore we propose the framework of security of biometric template, based on the concept of encryption with secret key, based on braid group for unauthorized user access, Fuzzy error correction code determine and encounters error if any error is introduce during the transmission between users to database server We store the biometric templates in encrypted form both without the fusion of score level and decision level in database server Stenography is related to hide biometric templates with help of secret key based on braid groups.

Section 2 of this paper describe the state of Biometric System, and section 3 as our proposed method and finally Section 4 the Conclusion

## Keywords:

Cryptography, Fuzzy Commitment Scheme, Biometric System Templates, Registration Phase, Verification Phase, Braid Groups

## **1.Introduction:**

In Ancient times the people recognized each other by face, voice and their tongs movement. Today's scenario is different due to large increase in our population. As a byproduct of Biometric system is activity in our day to day communication, we need more secure and authenticated communication. This requirement challenged us of think on biometric facilitation[5].Large number of government organizations, industry and non government organizations to must attain security with accuracy and error free Communication .The various properties exhibited by biometric system are uniqueness collectability, performance, acceptability and circumvention [3]

Various cryptography approaches are used for protection of biometric templates some are based on the hardware, some are software base, but common encryption technique AES and RAS cannot be used, because of large number of interclass variation in the biometric templates [5][6][16].Personal base authentication system use behavioral or physiological feature to identify one person form other. Many application for biometric system are available and all of them fall into two main functionality verification and identification [5][6][7].

A large number of errors in biometric templates are encountered in day to day communication due to environments change. Many approaches used to solve the noise problem which is one approach Fuzzy commitment scheme is one the scheme for solving the noise problem in biometric templates of a biometric recognition system, Jules and Wattenberg's Fuzzy Commitment Scheme[3]has been published to handling difference occurring

between two captured of biometric data, using error correcting code.

#### 1.1 Related work

A handful of papers have been published so far in the area of key generation for biometric secure templates [3,4,10,12,13] Further more amount of work suggest biometric data gathering one or more actual biometric analysis and combine their results which increase the reliability of the biometric system[1][2]

## **2.Preliminaries**

#### 2.1 Biometric System

Biometric System measure and analyzes biological data and provide capability of identifying person based on Their intrinsic characteristics that can be physical such hand shape, a fingerprint, facial characteristics, voice, or DNA and proving that the aim is same that is registered in the biometric security system as a biometric template Basically Biometric-Based Personal authentication systems have five major components 1.Sensor, 2.Feature-Extractor, 3.Template Database Server, 4.Matcher, 5.Decesion Module. Sensor is the interface between users and Biometric System and its function is to scan biometric traits of the user. Feature extractor module process the scan biometric traits to extract the silent feature set, in some case Feature. Extractor is followed by quality assessment module for checking the sufficient quality of scan biometric traits In Biometric-Based Personal authentication systems works in two steps

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**Step1**.Registered: In the first step the scan biometric traits of user R are processed and then stored in the database Server in Encrypted or secure form

**Step2.** Verification: In this step again new scan biometric traits of user R 'is captured and processed its then process for Comparing it with already Registered scan biometric traits as template and if both R and R' is matches then it is valid for Validation in biometric system otherwise Verification is invalid in biometric system Basically two types of errors are introduce In the Biometric-Based Personal authentication systems[14]

1. False Reject: Valid user (Probability of un matching of two traits R and R' of same user)

2. False Acceptance: Invalid User (Probability of matching of two traits R and R' of same user)

#### 2.2 Error Correction Code

**2.2.1 Definition:** A Metric space is a set B with a distance function dist B X  $B \rightarrow B^+=\{0, \infty)$ , which obeys the usual properties (symmetric, triangle inequalities, Zero distance between equal points) [13].

**2.2.2 Definition** : Let  $B(0,1)^n$  be a code set which consists of a set code words  $b_i$  of length n ,The distance metric between any two code words  $b_i$  and  $b_i$  in B is defined By

dist ( $b_i, b_j$ )= $\sum_{r=1}^{n} |b_{ir} - b_{jr}|$   $b_i, b_j \in \mathbb{C}$ This is known as Hamming Distance [13].

**2.2.2 Definition:** An Error Correction Function f for a code B Defined as

 $F(b_i) = \{b_j/\text{dist}(b_i, b_j) \text{ is the minimum, over } B - \{b_i\}\}.$ Here,  $b_{j=} f(b_i)$  is called the nearest neighbor of  $b_i[16]$ .

**2.2.3 Definition:** the measurement of nearness between two code word v and v' is defined by nearness (v,v')=dist(v,v')/n, it is obvious that  $0 \le nearness(v,v') \le 1[13]$ .

**2.2.4 Definition:** the fuzzy membership function for a code v' to be equal to given v is defined as[17]

FUZZ (v') =0 if nearness (v,v') = $z \le z_0 \le 1$ )

=z otherwise

2.3 Braid Group

Emil Artin[11] in 1925 defined B, the braid group of index n, using following generators and relations: Consider the generators  $\sigma_1$ ,  $\sigma_2$ ......,  $\sigma_n$  where  $\sigma_i$  represents the braid in which the(*i*+1)<sup>st</sup> string th

crosses over the i string while all other

Otherwise rejected strings remain uncrossed. The defining relations are

$$\sigma_i \sigma_{j=} \sigma_j \sigma_i / i - j \geq 2,$$
  
$$\sigma_i \sigma_j \sigma_{i=} \sigma_j \sigma_i \sigma_j \text{ For } |i-j| = 1$$

n-braid has the following geometric An interpretation: It is a set of disjoint n-strands all of which are attached to two horizontal bars at the top and at the bottom such that each strands always heads downward as one walks along the strand from the top to the bottom. In this geometric interpretation, each generator  $\sigma_{ii}$  represents the process of swapping the  $i^{th}$  strand with the next one (with  $i^{th}$  strand going under the  $(i+1)^{\text{th}}$  one). Two braids are equivalent if one can be deformed to the other continuously in the set of braids. Bn is the set of all equivalence classes of geometric n-braids with a natural group structure. The multiplication ab of two braids a and b is the braid obtained by positioning a on the top of b. The identity *e* is the braid consisting of *n* straight vertical strands and the inverse of *a* is the reflection of a with respect to a horizontal line. So  $\sigma^{-1}$  can be obtained from  $\Box$  by switching the overstrand and understrand. Δ=( $\sigma_1$ ,  $\sigma_2$ ,.....,  $\sigma_{n-1}$ ) ( $\sigma_1$ ,  $\sigma_2$ ,....,  $\sigma_{n-2}$ ) ( $\sigma_1$ ,  $\sigma_2$ )  $(\sigma_i)$  is called the fundamental braid. We describe some mathematically hard problems in braid groups. We say that x and y are conjugate if there is an element a such that  $y = \Box axa^{\Box f}$ . For m < n;  $B_m$  can be considered as a subgroup of  $B_n$  generated by  $\sigma_i$ , σ<sub>2</sub>,..... σ<sub>m-1</sub>

## 3. Proposed Scheme

Here we propose biometric based Authentication system using the concept of braid group based cryptosystem for security purpose and fuzzy commitment schema for error correction. We describe the complete process step of biometric Authentication system will achieve the high security and accuracy

**Step 1**: Scanning Scan the biometric traits of users using sensor at the user side 1. False Reject: Valid user (Probability of un matching of two traits R and R' of same user)

**Step2**: Quality Assessment Checking This process is to check the quality is sufficient of scanned traits of user for further processing

**Step3:** Process for convert scans biometric traits of users to Message Bits

First the scan biometric traits images decomposed into w/8 \*h/8 blocks where each one contain a fix number of pixel where w=height and h=height

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**Step 4:** Feature Extraction: At this step the silent feature are extracted from transformed block of biometric traits and generate the biometric templates as set of blocks

**Step 5:** Cryptographic key generation based on braid Group

Our scheme is made up of Five algorithms: setup, encryption, decryption, enrollment, verification

**Initially set up phase**: the environment is set up initially for generating public and private key,the secret key according to the algorithms using braid group at time  $t_o$  where event time is function  $E=\{t_i,a_i\}$  of time and event algorithms and publish secret key between two parties requires for secures communication starting

**Encryption**: According to the commit phase, Bob commits to a Message  $m \in M$  to Alice, send as:

B<sub>n</sub>:Sufficient complicated Braids R. Random Braids Bob Alice Private key: Bobprivate where Bobprivate€LEn Alicencivate Where Alicencivate€UEn Public key : Bobpublic = Bobneivate Bn Bob-1 private Alicenthic = Alicentivate Bn Alice-1 private Identity: ID<sub>Bob</sub> where ID<sub>Bob</sub>€L<sub>Rn</sub> ID<sub>Alice</sub>where ID<sub>Alice</sub>€U<sub>Rn</sub> Calculation: CEab=Bobnublic(IDEab)@Bobneivate(IDEab) CAlica=Alicenublic(IDAice)0 Alicencivate (IDAlice) Alice Decrypt Bobneivate (IDBob) with Alice Public Bob Decrypt Alicensivate (IDAlice) with Alice Public Key Alicenthid And find the IDAlice and calculate Key Bobpublic And find the ID<sub>Bob</sub> and calculate CBob@ Bobnrivan(IDRob)= Bobnublic(IDRob) and then CAlice? Alicencivate (IDAlice)= Alicenublic (IDAice) and then encrypt encrypt Its with  $Boh_{nublic}$  and find the  $ID_{Alice}$  is know as Identity of Alice Its with Alicenublic and find the IDBob is know as Identity of Bob Firstly, to find out the identity of each other, both make some calculation as Bobmag= Bobmivate Bn Bob-1 private Dicekey = Alicentivate Bobnublic Alice<sup>-1</sup>private And Alice Calculate = AliceKev(IDAice)Bobmar(IDAlice)<sup>-1</sup> AliceKev<sup>-1</sup> Alicense Bobkey = Bobneivate Alicenublic Bob-1 private And Bob Calculate BobKey=(IDAice) (BobKey) Alicemse( BobKey<sup>-1</sup>)( IDAlice)-1 In each above step, if Bobker = Aliceker or is I, then the protocol run is terminated with failure. So both Alice and Bob have secret Kchared = bobker = Aliceker and both can communicate secretly for that session.

c=commita lg(XOR,g(m1),  $F_{kshared}$  (m),R) where  $F_{kshared}$  (m)=m<sub>1</sub>K<sub>shared</sub>+e, m is the K-bit message , m=m<sub>1</sub>\Delta m<sub>2</sub>,  $F_{kshared}$  (m) is an n-bit unreadable text ,R is random braid (R can be changed in every Message that sent),e=h(m2),here h is an invertible function which maps m<sub>2</sub> in to an n-bit error vector of weight  $\alpha$ 

**Decryption**: in the open Phase Alice open the commitment at time  $t_2$  which was sent by Bob at time  $t_1$ using the inverse procedure of commita  $lg(a_2)$ and Alice make a calculation c' using with help of secret key c'=commita  $lg(XOR,g(m1), F_{kshared} (m),R)$ and check its time t3 result is same as the sent by Bob Fuzzy decision Making

If  $(near(c,c') \leq Zo)$  if result is less than Accept otherwise reject

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Enrollment/Identification Phase of Biometric System

#### In the Identification phase

Biometric template's of user is u' at receive at receive side and compared it with one is stored previously in database if they are matching then user is validate for the system

#### In the Registration Phase

User biometric scan traits as biometric template as encrypted using secret key and send to

Database server where database server found encrypted biometric templates with error are introduce in transmission channel so finally  $E_v(M)=Ma+@e$ 

#### Step 6 error correction codes

If Error introduce in transmission of biometric template we can correct using fuzzy error correcting code if any error are introduce during the transmission of biometric templates

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Receiver check the dist(t(c),c')>0 he realize that there is an error occur during transmission .receive applies the error correction function c':f(c)

The receive will compute nearness  $(t{c},f{c}=dist{t{c},f(c')}/n$ Fuzz(c')=0 if nearness(c,c')=Z<Z<1 =z otherwise

## 4. Conclusion

This paper presents fuzzy commitment Method for error correcting code and braid group to compute a cryptographic key for biometric data .The method not only meet the security requirements, but also produce comparable accuracy to well-known K-NN classification method. Experiments on real biometric data, particular fingerprints and voice, are being conducted and will be reported in the near future.

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